

Performance Analysis of Cooperative Beacon Sensing Strategies for Spatially Random Cognitive Users

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AQ1 1 **Abstract**—Primary user (PU) beacons must be detected by
 2 cognitive users (CUs) to access spectrum holes, and misdetection
 3 results in interference on PUs. To alleviate this problem, sensing
 4 results of spatially separated CUs can be combined to make a
 AQ2 5 final decision. In this paper, we analyze several such cooperative
 6 beacon sensing (CBS) strategies given spatial randomness of CU
 7 and PU nodes, which is modeled via independent homogeneous
 8 Poisson point processes. We consider two cases of beacon emitter
 9 placement: 1) at PU-transmitters and 2) at PU-receivers. We
 10 analyze three separate local beacon detection schemes and propose
 11 five CBS schemes. They require the sharing of CU results
 12 via a control channel subject to Rayleigh fading and path loss,
 13 and making a final decision via the OR rule. By using stochastic
 14 geometry, we derive both the misdetection probability, the
 15 false alarm probability, and the primary outage and show that
 16 impressive gains are achievable. For example, with PU-receiver
 17 beacons, CBS reduces misdetection by a factor of 10^4 . In contrast,
 18 with PU-transmitter beacons, the reduction diminishes with
 19 the increased cell radii, but there exists an optimum cooperation
 20 radius.

AQ3 21 **Index Terms**—

22 I. INTRODUCTION

23 **T**HE MISDETECTION of beacon signals emitted by pri-
 24 mary users (PUs) by cognitive users (CUs) is a major
 25 problem, leading to interference on PU nodes which reduces
 26 their data throughput and increases their outage. Thus, fixing
 27 the beacon misdetection problem is critical to the deployment
 28 of cognitive radio (CR) networks. The CR paradigm is driven
 29 by the scarcity of spectrum and its inefficient use, two of the
 30 most critical challenges facing modern wireless networks [2].
 31 For example, traditional static spectrum assignments to indi-
 32 vidual users/services lead to 85% or more idle licensed
 33 spectrum [3]. Thus, unlicensed (i.e., cognitive) opportunistic
 34 access to licensed spectrum [4] has been standardized in IEEE
 35 802.22 Wireless Regional Area Network (WRAN) and its
 36 amendments, IEEE 802.11af for wireless local area networks,

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37 licensed shared access (LSA) for Long Term Evolution (LTE)
 38 and others [5]. In particular, the cognitive interweave mode
 39 aims to allow opportunistic access to temporary unused
 40 space-time-frequency slots (spectrum holes) [6]. However, CU
 41 devices must then accurately detect active PU transmissions
 42 in real time via matched filtering, cyclostationarity, energy,
 43 eigenvalues, beacons or other methods [7]–[10].

44 Of these, PU beacon signaling has the benefits of effi-
 45 ciency and simplicity [11]–[16]. Grant or denial beacons
 46 are simply out-of-band, on-off modulated electromagnetic
 47 waves [17], proposed for IEEE 802.22.1 [18] and cognitive
 48 cellular systems [17], [19]. In this work, we focus on the
 49 problem of detecting **denial beacons** of active PU nodes.
 50 Beacon misdetection, which leads to interference on the PUs,
 51 occurs due to multipath fading, path loss, receiver uncer-
 52 tainty and other factors [20], [21]. Thus, a classical solution
 53 is to exploit spatial diversity. We can thus use multiple bea-
 54 con measurements from spatially separated CUs and combine
 55 them into one final decision. This is an instance of cooper-
 56 ative sensing, which can be based on OR, AND, or majority
 57 rules [8], [22]. In this paper, we will limit ourselves to the OR
 58 rule to determine the presence of a denial beacon, which leads
 59 to conservative spectrum access attempts (i.e., ensuring less
 60 interference). The reduction in misdetection probability due
 61 to cooperative beacon sensing (CBS) depends on the number
 62 of cooperating CUs and their locations [23], which are ran-
 63 dom. Due to this spatial randomness, path loss, and fading, the
 64 expected performance improvements of CBS may be severely
 65 compromised. To characterize such issues, a comprehensive
 66 analysis of the overall beacon misdetection probability (P_{md})
 67 is necessary.

68 A. Problem Statement and Contribution

69 In this paper, we analyze the overall P_{md} and false alarm
 70 probability (P_f) of several CBS methods as a function of how
 71 cooperating CUs are selected, local detection methods, spatial
 72 randomness of primary and secondary nodes, channel fading,
 73 and the sharing of imperfect decisions. Specifically, we address
 74 the following questions: 1) How does a CU device locally
 75 process one or more beacons transmitted from multiple PU
 76 devices to mitigate the impact of fading and path loss? 2) How
 77 do we select a set of CUs for cooperative spectrum sensing
 78 when the beacons are sent by PU-receiver nodes or PU-
 79 transmitter nodes? What are the rules that specify a suitable

80 set of cooperating CUs? The cooperative sensing phase will be
 81 affected by the channel propagation characteristics and spatial
 82 randomness of the cooperating CUs. The availability of chan-
 83 nel state information (CSI) for the CU-to-CU channels affects
 84 the selection of best nodes to cooperate with. Clearly, the coop-
 85 erating set should be chosen to minimize P_{md} , which will
 86 depend on mutual distances and fading conditions. 3) What is
 87 the overall performance of CBS?

88 To investigate all these questions for coexisting cellular (pri-
 89 mary) and cognitive networks, we first ensure that the spatial
 90 randomness of nodes is fully accounted for. To this end, we
 91 use the tools from spatial geometry to model the random loca-
 92 tions of PU and CU nodes. Specifically, we model PU-receiver
 93 nodes and CUs as Poisson Point Processes (PPPs) [24].
 94 However, the PU-transmitters are fixed at the centers of hexag-
 95 onal cells. For realistic propagation modeling, we incorporate
 96 both power-law path loss and Rayleigh fading. The beacon
 97 detection process of a CU is consisted of two distinct phases:
 98 1) local detection, and 2) cooperation. The sharing of detec-
 99 tion results is done via a control channel subject to fading
 100 and path-loss. Moreover, we consider beacons sent by PU-
 101 receivers (Case 1) and by PU-transmitters (Case 2). Our main
 102 contributions in this paper are as follows:

- 103 i) For phase one, we propose three local beacon processing
 104 schemes: 1) aggregating beacon powers, 2) separately
 105 sensing multiple beacons, and 3) detecting the best
 106 average received beacon signal (i.e., from the closest).
- 107 ii) For phase two, we propose three cooperation schemes:
 108 1) nearest scheme, 2) multiple-random scheme, and
 109 3) best received power scheme. For beacons emitted
 110 by PU-transmitters, we propose two additional schemes:
 111 1) nearest CU to PU-transmitter scheme and 2) random
 112 CU to PU-transmitter scheme.
- 113 iii) For all these schemes, we derive P_{md} and P_f from the
 114 OR rule fusion in order to characterize the performance
 115 improvement of CBS under different system parameters.
- 116 iv) We derive the outage probability of a PU-receiver
 117 to characterize how its performance is affected by
 118 interference due to beacon misdetection.

119 B. Prior Research

120 We first review papers that do not focus on beacons signal-
 121 ing but perform general misdetection analysis and interference
 122 characterization for CR networks [15], [25]–[30]. For brevity,
 123 we denote the aggregate interference by I . In [15], the distri-
 124 bution of I is characterized in terms of sensitivity, transmit
 125 power, density of the CUs, the propagation characteristics,
 126 and cooperative spectrum sensing. In [30], the theory of trun-
 127 cated stable distributions and power control are studied for a
 128 CR network. Reference [25] analyzes the primary coverage
 129 probability under misdetections and false alarms, and devel-
 130 ops an approximation and bounds for the Laplace transform of
 131 I . Statistics of I from a secondary network with an ALOHA
 132 based medium access control, spectrum sensing, and power
 133 control is derived [26]. Moreover, [27] derives the moment
 134 generating function of I for a spectrum sensing CR network,
 135 and a scheme is proposed to maximize the transmission powers

of multiple active CU transmitters while satisfying I con- 136
 straints. This scheme leads to significantly higher capacity. 137
 Reference [29] analyzes the geometric region allowing CR 138
 transmission with the help of cooperative sensors, and finds 139
 that the shape of this region is not circular. Furthermore, [31] 140
 develops models for bounding interference levels by modeling 141
 CUs as a modified Matern process. Co-operating spectrum 142
 sensing methods are analyzed over correlated shadow fading 143
 environments [28]. The spatial throughput of a CR network is 144
 characterized for a two threshold based opportunistic spectrum 145
 access protocol in [13]. 146

Several works consider spectrum sensing using 147
 beacon detection and also cooperative spectrum sens- 148
 ing [11], [13], [32]–[35]. Reference [11] analyzes 149
 capacity-outage probability of a PU due to interference 150
 from beacon misdetection. The emission of beacons by PU- 151
 receiver nodes leads to higher capacity-outage performance. 152
 Furthermore, [34] considers three levels of cooperation under 153
 beacon transmissions from the primary users. It is shown 154
 that cooperation is vital when the CU node density is high. 155
 Threshold based opportunistic spectrum access methods 156
 are studied in [13] under PU-transmitter and receiver pilot 157
 signals and beacons, and the spatial opportunity (probability 158
 that an arbitrary location is discovered as a spectrum hole) 159
 is derived. Furthermore, [32] and [33] study the resultant 160
 aggregate interference due to misdetection in beacon based 161
 CR networks. Moreover, [35] studies the soft combination 162
 of spectrum information shared by the cooperating nodes 163
 when for multiple beacon signalling, and derives the optimal 164
 beacon sequence to reduce misdetection. 165

The differences among the aforementioned works and this 166
 paper are now described. First, spatial randomness of CUs 167
 is not considered in [11] and thus the spatial densities of 168
 the nodes do not appear in their analysis. Second, the exist- 169
 ence of multiple PU-receivers is not considered [32], [33]. 170
 Third, the control channel for sharing the sensing result 171
 has been assumed perfect [11], [13], [34]. In contrast, in 172
 this paper consider the effect of propagation impairments 173
 (path loss and fading) on the quality of reception of control 174
 signals. Fourth, the availability of channel state informa- 175
 tion (CSI) has not been considered for cooperating node 176
 selection [11], [13], [32]–[35]. However, we CBS strategies 177
 depending on the availability of CSI. Fifth, no distinction 178
 is made between beacons emitted by PU-transmitters and 179
 those by PU-receivers [32], [33]. In contrast, this paper 180
 derives the interference statistics of the two cases in detail. 181
 Sixth, the impact of spatial locations has not been consid- 182
 ered [11], [13], [34], [35]. As such, our paper strives to fill 183
 these gaps while investigating the misdetection probability 184
 reduction of cooperative sensing. 185

This paper is organized as follows. Section II introduces 186
 the signal model including the spatial model, signal propaga- 187
 tion, local detection schemes, and cooperation schemes. The 188
 misdetection probability P_{md} is analyzed for PU-receiver and 189
 PU-transmitter beacons in Sections III and IV. Section V char- 190
 acterizes the primary system performance. Numerical results 191
 are provided in Section VI while Section VII concludes the 192
 paper. 193

TABLE I
LIST OF COMMONLY USED PDFS

Name	PDF
$Lin(\alpha)$	$f(t) = \frac{2t}{\alpha^2}, 0 < t < \alpha$
$Ral(\alpha)$	$f(t) = 2\alpha t e^{-\alpha t^2}, 0 < t < \infty$
$TRal(\alpha, \beta)$	$f(t) = \frac{2\alpha t e^{-\alpha t^2}}{1 - e^{-\alpha \beta^2}}, 0 < t < \beta$

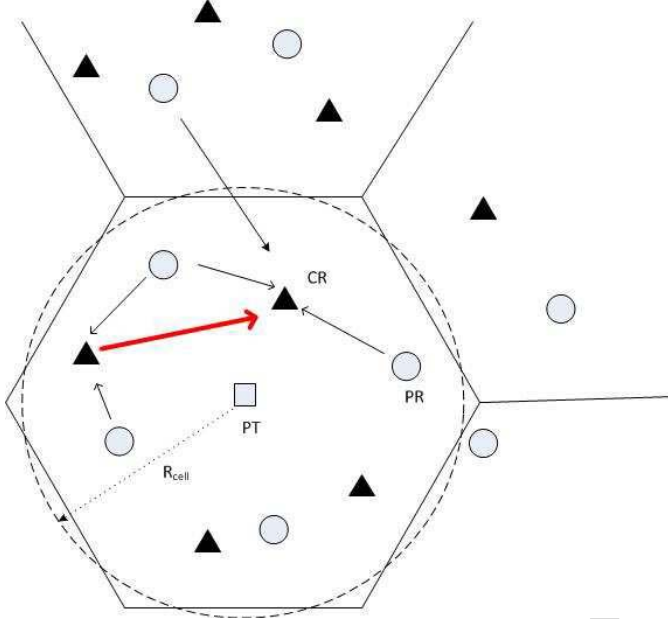


Fig. 1. PU-receiver node emit beacons. Squares, triangles, circles, and solid arrows respectively denote the PU-transmitters, CUs, PU-receivers, and the beacon signals. Each cell is hexagonal with a PU-transmitter at the center. PU-receivers and CUs are distributed randomly in \mathbb{R}^2 .

194 *Notations:* $\Gamma(w, a) = \int_a^\infty t^{w-1} e^{-t} dt$ and $\Gamma(w) =$
 195 $\Gamma(w, 0)$ [36]. $\Pr[A]$ is the probability of event A , $f(\cdot)$ and $F(\cdot)$
 196 are the probability density function (PDF) and the cumulative
 197 distribution function (CDF), $M_X(\cdot)$ is the moment generating
 198 function (MGF), $E_n(\cdot)$ is the generalized exponential integral,
 199 and $E_X[\cdot]$ denotes the expectation over random variable X . The
 200 Euclidean distance between two points x and y is denoted by
 201 $\|x - y\|$. The following PDFs (Linear, Rayleigh, and truncated
 202 Rayleigh) listed in Table I will be used commonly throughout
 203 the paper.

204 II. SYSTEM MODEL

205 A. Spatial Distribution

206 We consider coexisting primary and cognitive (secondary)
 207 networks. We assume the PU network to be of a conventional
 208 cellular type with different cells using the same frequency
 209 set (the frequency reuse factor is 1). The area is divided into
 210 hexagonal cells with a PU-transmitter (e.g., base-station) at the
 211 center of each (Fig. 1), which serves a set of spatially random
 212 PU-receivers within each cell. The cognitive network which
 213 can be an ad-hoc network or a sensor network [37] utilizes
 214 primary spectrum holes to transmit data. To facilitate anal-
 215 ysis, we approximate the hexagonal cells with circular cells
 216 having a radius of r_{cell} (Fig. 1). The spatial randomness of
 217 CUs is also considered.

218 To model spatial randomness, we will make use of point
 219 processes. For our purposes, a point process Φ is a collection
 220 of points $\{x_1, x_2, \dots\}$ where $x_k \in \mathbb{R}^2$ is a point represent-
 221 ing the location of a radio node. We say Φ is a Poisson
 222 point process with rate $\lambda > 0$ if (1) the number of nodes
 223 within a bounded area A denoted by $N(A)$ is a Poisson ran-
 224 dom variable with $\mathbb{E}[N(A)] = \lambda A$ and (2) the number of
 225 nodes in two non-overlapping areas are independently dis-
 226 tributed [38]. Poisson processes are widely used to model the
 227 locations wireless nodes due to their mathematical tractability
 228 and accuracy [30], [39].

229 In this paper, we model PU-receivers and CU nodes as two
 230 independent homogeneous PPPs Φ_p and Φ_s in \mathbb{R}^2 with spatial
 231 densities λ_p and λ_s . Thus, the number of nodes within the
 232 bounded area A is given by

$$233 \Pr[N(A) = n] = \frac{(\lambda A)^n}{n!} e^{-\lambda A}, n = 0, 1, 2, \dots \quad (1)$$

234 where $\lambda \in \{\lambda_p, \lambda_s\}$ [38].

235 We assume that the CSI of the PU-CU links are not available
 236 to individual CUs. This assumption is reasonable and com-
 237 mon [11] because of the general commercial and regulatory
 238 pressures that push primary and secondary networks to operate
 239 independently. However, a CU may or may not know about the
 240 CSI of links between itself and other CUs. The degree of
 241 availability of this CSI to CUs will impact the development
 242 of cooperative spectrum sensing protocols.

243 In this work, mobility of wireless nodes is not analyzed
 244 for two reasons. First, some PU nodes are fixed (e.g., base-
 245 stations, TV receivers and others). Second, even if the CUs
 246 move randomly (e.g., random walk or the Brownian motion),
 247 a snapshot of at any specific time generates a homogeneous
 248 PPP. Nevertheless, the impact of the mobility of nodes is a
 249 challenging, future topic.

250 Furthermore, we assume that CUs are always ready to trans-
 251 mit data upon detecting a spectrum hole and that all the
 252 PU-receivers are active. There is no loss of generality in these
 253 assumptions since activity factors (≤ 1) can easily be incor-
 254 porated using the Coloring Theorem [24]. That is, if nodes of
 255 a PPP Φ with intensity λ are marked independently, and p_t
 256 is the probability of a node receiving the t -th color, the set of
 257 t -th color nodes forms a PPP Φ_t with intensity $p_t \lambda$. Thus, if a
 258 PU-receiver is active with an activity factor of q_p , the set of
 259 active PU-receivers follows a thinned PPP with intensity $q_p \lambda_p$.
 260 The same argument holds for the CUs.

261 B. Signal Propagation

262 The propagation effects are characterized by independent
 263 Rayleigh fading and log-distance path loss [40]. With small-
 264 scale Rayleigh fading, the channel power gain $|h|^2$ has the
 265 Exponential PDF $f_{|h|^2}(t) = e^{-t}, 0 < t < \infty$. The log-
 266 distance path loss model specifies that the received power
 267 $P_R = P r^{-\alpha}$ where r is the distance between the transmitter
 268 and the receiver, P is the transmit power and α is the path
 269 loss exponent. The path loss exponent is a function of carrier
 270 frequency, terrain, obstructions, antenna heights and others.
 271 The typical values range from 2 to 8 (at around 1 GHz).
 272 Note however that because $g(r) = r^{-\alpha}$ leads to analytical

difficulties when $r \rightarrow 0$, we will also use $g(r) = \min(1, r^{-\alpha})$. Both forms of $g(r)$ will yield the same results because since spatial densities are small (e.g., $\lambda_p, \lambda_s \ll 1$), the probability that the distance is small is negligible, $P[r < 1] \rightarrow 0$.

Throughout the paper, we assume that all CUs transmit at a fixed power level [41]–[43]. Although CU power control methods are beyond the scope of this paper, they can be easily incorporated if needed [44].

C. Local Detection

As mentioned before, beacon detection process at a CU is divided into 2 phases: the local detection phase, and the cooperative phase. In this paper, we assume the downlink transmission of the cellular network with denial beacons where the PU devices (either PU-transmitters or PU-receivers [16], [17]) transmit a beacon signal. This beacon will have a set number of bits indicating that κ ($\kappa \in (1 \dots K)$) future time-slots will be occupied by the transmitting device. Moreover, the beacon would uniquely identify the transmitting PU device, and would enable synchronization between the primary and secondary network. Furthermore, the beacon signal is transmitted before channel access by the PU device. For example, in the case of PU-transmitter beacons, the device sends the beacon signal before transmitting its data, while for PU-receiver beacons, the beacon is emitted by all active devices before they begin receiving oncoming data.

Beacons emitted by PU-receivers are more likely to be correctly heard by CUs which can interfere the most. However, PU-receivers (e.g., hand-held user devices), will increase their battery drain because of beacons emissions. To counteract this, beacon signals can be made shorter, their frequency can be reduced, or their power can be reduced. All these options may unfortunately increase the miss detection of beacons. On the other hand, when PU-transmitters emit beacons, the CUs which can potentially interfere the most for cell-edge PU-receivers may miss them. Nevertheless, such beacons can be used under high PU-transmitter densities (lower cell radii), and where PU-receivers are severely power limited [11]. Otherwise, PU-receiver beacons should be used wherever possible.

1) *PU-Receiver Beacons*: Without the loss of generality, we assume that all PU-receivers are active and transmit beacons. For this to work, we assume synchronization between the different PU-receivers. However, if only a subset of the PU-receiver nodes are active, this can be easily incorporated using the Coloring Theorem [24]. Note that a CU may detect a beacon from a PU-receiver in another cell (Fig. 1). Thus, we suggest three local beacon detection schemes. These schemes are:

i) *Aggregating all beacons in the range*: Each CU simply uses the aggregate beacon power received, which does not require it to differentiate among the different PU-receiver beacons. However, this is a conservative approach in terms of opportunistic spectrum access because the aggregate beacon power may exceed the sensing threshold even when nearby PU-receivers are inactive.

ii) *Sensing beacons separately and OR combining them*: A CU is assumed to differentiate the beacons emitted by various PU-receivers (e.g., each one may use a different orthogonal code [45] or matched filtering may be used [13]). Thus, each distinct beacon is uniquely sensed. However, the implementation of a separate beacon sensing scheme has significant challenges. As the spatial density of PU nodes increases, this schemes requires additional processing. Moreover, longer code-words and thus longer beacons are needed to uniquely identify the different PU-receivers. On a practical point of view, only the PU-receivers within a certain radius from the CU may be considered for local detection instead of all the PU-receivers within the geographical area. The separate sensing scheme is advantageous for CUs because it allows them to access the spectrum whenever a beacon signal from a PU is less than the threshold. This is in contrast with the aggregate scheme where even if the individual beacon powers are far less than the threshold, the aggregate can still be above the threshold, barring a CU from accessing the spectrum.

iii) *Sensing the beacon from the closest PU-receiver only*: The CU must find the closest PU-receiver perhaps by measuring the average received signal power [46]. Moreover, the CU must differentiate among the beacons from different PU-receivers in order to achieve this. This scheme has the advantage of considerable less processing than the separately sensing scheme after the closest PU-receiver has been established. Moreover, it provides the best opportunities for a CU to access the spectrum among the three local detection schemes. However, because only a single PU beacon is considered, there is a high misdetection probability.

2) *PU-Transmitter Beacons*: We assume that all PU-transmitters become active at the same time. Each CU listens to its own cell's PU-transmitter for beacon signals. It should be noted that while a CU may receive a better instantaneous signal from a neighbouring cell due to a favourable channel, the PU-transmitter of its cell would also be the closest PU-transmitter to a given CU, and thus would provide the best received beacon signal power on average. We assume that the CUs have the ability to uniquely identify its own PU-transmitter from neighbouring PU-transmitters.¹ While beacon signal reception from out-of-cell PU-transmitters can also be considered, we leave this for future work.

D. Co-Operative Sensing

In the cooperative phase, the CU will select one or more other CUs to obtain the sensing results via a single narrow-band control signal. We assume that the CUs can identify each other via the use of separate orthogonal codes or time slots. In our analysis, we will consider distributed cooperation schemes without the involvement of a fusion center, information sharing via decision-fusion, and combination via the OR rule [8]. The OR rule minimizes P_{md} compared to other combining

¹Separately identifying PU-transmitter beacons may be achieved by using unique codes or time slots.

rules [8]. Because distributed co-operating schemes are used, each individual CU keeps a dynamic database of neighbouring CUs. This database will include details about activity, distance, and CSI if available. Information for the individual databases is obtained via periodic control signals, and updated regularly. We thus propose three cooperation schemes, where the selection is based on the information within each CU's database. They are:

- i) *Nearest scheme*: Each CU cooperates with its closest neighbor CU, which provides the best received signal power on average. To implement this, distances among the CUs are needed [47]. These distances may be obtained via a database, shared GPS information or via periodic control signals.
- ii) *Multiple random scheme*: Here, M neighbouring CUs are randomly selected within a cooperation radius of R_c ($\ll R_e$). A CU is assumed to only cooperate with a neighbour within this radius. The signals from nodes beyond the outer distance R_e are assumed to have negligible power due to high path loss. If the number of CUs within R_c is less than M , all would be selected. The selected nodes are always available for cooperation.
- iii) *Best received power scheme*: In this scheme, each CU cooperates with the neighbouring CU providing the best instantaneous received signal power. This amounts to the lowest propagation loss considering both path loss and fading. We assume that each CU knows CSI and the positions of other CUs. Moreover, we further assume that a CU can cooperate with nodes outside its own cell.

We will assume that CUs can differentiate the beacon signals from the PU-receivers and the control signals from other cooperating CUs. For example, this involves using separate orthogonal codes for different CUs and PU-receivers, using different time slots, matched filtering, or having a separate narrow band channel for CU spectrum information sharing [13], [14], [45]. Furthermore, it should be noted that each CU shares its local detection result, but not the final decision of CBS.

With PU-transmitter beacons, we propose two additional schemes based on the intuition that CUs close to the PU-transmitter will have a better chance of correctly detecting the beacon. These schemes are:

- i) *Nearest CU to PU-transmitter scheme*: Each CU, $x \in \Phi_s$, selects the closest CU to the PU-transmitter, which has the best probability to detect the beacon signal due to the lowest path loss. Furthermore, selection of distances to a fixed PU-transmitter may be less complex than find all CU-to-CU distance.
- ii) *Random CU to PU-transmitter scheme*: A random CU within a distance of R_c from the PU-transmitter is selected. The distance constraint from the PU-transmitter which ensures the cooperating CU has a good chance of detecting the PU beacon. This scheme has the advantage over the previous scheme of not burdening a single CU (the one closest to the PU-transmitter) for sensing data.

Choosing other CU nodes to cooperate with based on distances to PU nodes is most suitable when PU-transmitters emit beacons. PU-transmitters would generally be fixed, and their

locations would thus not change dynamically. As such, choosing CU nodes within a certain distance from the PU-transmitter is relatively straightforward. On the other hand, PU-receivers may be fluid in their activity, and multiple PU-receivers will be transmitting (with PU-transmitters, we assume the CU only listens to the PU-transmitter of its own cell) their beacons. As such, choosing cooperating CU nodes satisfying distance requirements from PU-receivers is more cumbersome, and such schemes are not considered in this paper.

III. P_{md} ANALYSIS FOR PU-RECEIVER BEACONS

A. Local Primary Beacon Detection

In this section, we analyze P_{md} for the local spectrum sensing methods in Section II-C.

1) *Aggregating Beacon Power*: Consider the CU node $x \in \Phi_s$ and the PU-receiver node $y \in \Phi_p$. The distance between them is $\|x - y\|$. However, as this distance becomes large, $g(\|x - y\|) \rightarrow 0$. As such, the beacons emitted by PU-receiver nodes y such that $\|x - y\| > R_e$ are considered to be negligible, where R_e is an outer distance. Since x and y are two random points from two independent PPPs, we need the distribution of the distance $\|x - y\|$. However, because a homogeneous Poisson process is considered for Φ_p , its points are distributed randomly. Moreover, due to the outer distance, the area of node distribution is annular. Therefore, the CDF of $\|x - y\|$ can be obtained as [43]

$$F_{\|x-y\|}(t) = \frac{t^2}{R_e^2}, \quad 0 < t < R_e. \quad (2)$$

Thus, $\|x - y\|$ is distributed with PDF $Lin(R_e)$.

All PU-receiver nodes $y \in \Phi_p$ transmit a beacon signal of constant power level P_b . As the CU will aggregate these beacons, the received beacon power at CU x is given by

$$P_R = P_b \sum_{y \in \Phi_p} |h_{x,y}|^2 g(\|x - y\|), \quad (3)$$

where $h_{x,y}$ is the channel between nodes x and y , and this incorporates both path loss and small scale fading. The received signal to noise ratio (SNR) γ at CU $x \in \Phi_s$ becomes $\gamma = \frac{P_R}{\sigma_b^2}$, where σ_b^2 is the additive noise variance. The CUs can employ energy detection of the beacon channel or use a received power threshold. However, as shown in [11], even an energy detection based scheme can be approximated as a simple received power threshold based scheme with an appropriate threshold. Therefore, in our analysis, a beacon is detected whenever the received beacon power $P_R > P_{th}$, where P_{th} is the reception threshold.

Let $P_{md}(x)$ be the probability of PU beacon misdetection by the CU $x \in \Phi_s$ in its local-detection phase. This probability is given by

$$P_{md}(x) = \Pr[P_R < P_{th}] = F_{P_R}(P_{th}),$$

which is the CDF of P_R . This can be evaluated using an MGF based approach [41], [48]–[50]. Let $M_{P_R}(s)$ be the MGF of the received beacon power at $x \in \Phi_s$, which is defined as $M_{P_R}(s) = E[e^{-sP_R}]$. If $M_{P_{R,y}}(s)$ is the MGF of the received

491 beacon power from $y \in \Phi_p$, and N is a Poisson random
492 variable with mean $\pi R_e^2 \lambda_p$, we can write $M_{P_R}(s)$ as [41], [43]

$$493 \quad M_{P_R}(s) = E_N \left[(M_{P_{R,y}}(s))^N \right] = e^{\pi R_e^2 \lambda_p (M_{P_{R,y}}(s) - 1)}. \quad (4)$$

494 $M_{P_{R,y}}(s)$ is obtained as follows.

$$495 \quad M_{P_{R,y}}(s) = E \left[e^{-s P_b |h_{x,y}|^2 g(\|x-y\|)} \right] \\ 496 \quad = \int_0^1 \frac{1}{1 + s P_b R_e^2} \frac{2t}{R_e^2} dt + \int_1^{R_e} \frac{1}{1 + s P_b t^{-\alpha}} \frac{2t}{R_e^2} dt. \quad (5)$$

497 A closed-form expression for the second integral (5) appears
498 intractable. However, using the expansion $(1+t)^{-1} =$
499 $\sum_{k=0}^{\infty} (-t)^k$, $|t| < 1$, we derive a simplified expression as

$$500 \quad M_{P_{R,y}}(s) = \frac{1}{R_e^2} \left(\frac{1}{1 + s P_b} + \sum_{l=0}^{\infty} 2(-s P_b)^l \frac{R_e^{2-\alpha l} - 1}{2 - \alpha l} \right). \quad (6)$$

501 $F_{P_R}(t)$ can be obtained through the inverse Laplace transform
502 by $F_{P_R}(t) = \mathcal{L}^{-1} \left(\frac{M_{P_R}(s)}{s} \right)$, and replacing t with P_{th} gives
503 $P_{md}(x)$, $x \in \Phi_s$. Note that because a closed-form solution is
504 not apparent for $P_{md}(x)$, where $x \in \Phi_s$, numerical techniques
505 and approximations must be used.

506 Although aggregating beacon power decreases P_{md} , viable
507 spectrum access opportunities are also lost due to detect-
508 ing aggregated beacons even when there may not be any
509 PU-receivers close by to be hindered by interference.

510 2) *Separately Sensing Primary Beacons*: Misdetection
511 occurs only when all beacon sensing outputs fall below the
512 threshold. Thus we have $P_{md}(x) = (\Pr[P_{R,y} < P_{th}])^N$,
513 where $x \in \Phi_s$ and $P_{R,y}$ is the beacon power from $y \in \Phi_p$
514 received at $x \in \Phi_s$, and N is a Poisson random variable with
515 $\mathbb{E}[N] = \lambda_p \pi R_e^2$. The misdetection of the beacon from $y \in \Phi_p$
516 may be written as $\Pr[P_{R,y} < P_{th}] = E_{(\|x-y\|)} \left[1 - e^{-\frac{P_{th}}{P_b g(\|x-y\|)}} \right]$.
517 Thus, denoting $\|x-y\| = t$, the local misdetection probability
518 may be expressed as

$$519 \quad P_{md}(x) = e^{-\pi R_e^2 \lambda_p \left(\frac{e^{-\frac{P_{th}}{P_b}}}{R_e^2} + \frac{2}{R_e^2} \int_1^{R_e} e^{-\frac{P_{th}}{P_b t^{-\alpha}}} dt \right)}. \quad (7)$$

520 Because a closed-form solution for (7) appears impossi-
521 ble, we numerically evaluate this. A series summation based
522 simplification can be used to simplify (7) which results in

$$523 \quad P_{md}(x) = e^{-\pi R_e^2 \lambda_p \left(\frac{e^{-\frac{P_{th}}{P_b}}}{R_e^2} + \frac{2}{R_e^2} \sum_{k=0}^{\infty} \frac{\left(-\frac{P_{th}}{P_b} \right)^k}{k!} \left(\frac{R_e^{2+\alpha k} - 1}{2 + \alpha k} \right) \right)}. \quad (8)$$

524 However, more resources are required for separate sensing,
525 and is invariably more complex. Furthermore, the PU-receivers
526 need to be co-ordinated to send separately identifiable beacons.
527 This may not be practical for certain PU-receiver types such
528 as digital terrestrial television subscribers.

529 3) *Closest PU-Receiver Selection*: Each CU, $x \in \Phi_s$,
530 senses the beacon emitted by the closest PU-receiver. The clos-
531 est PU-receiver may be found in practice by measuring the
532 average received signal power [46]. Moreover, the CU must
533 then have the ability to differentiate among different beacons.

534 Let $y^* = \arg \min_{y \in \Phi_p} \|y - x\|$ ($y^* \in \Phi_p$) be the nearest
535 PU-receiver to $x \in \Phi_s$, and the distance $r^* = \|y^* - x\|$.

The distribution of r^* is derived via the void probability of
536 a PPP (probability of no nodes within a given radius from the
537 origin) [51], [52], and is found out to be $Ral(\pi \lambda_p)$.
538

539 However, as the beacons from node $y \in \Phi_p$ at a distance
540 more than R_e are neglected due to path loss, there may be
541 an occasion where there is no closet PU-receiver within R_e .
542 The probability of this event is $p_0 = e^{-\pi \lambda_p R_e^2}$. Whenever this
543 occurs, the CU $x \in \Phi_s$ will misdetect with probability 1.
544 However, conversely, because of the high path loss in such
545 a scenario, the interfering signals will also have a negligible
546 effect on the primary system. Let r_1^* be the truncated dis-
547 tance from x to y^* whenever $r^* < R_e$. Thus, r_1^* is distributed
548 according to $TRal(\pi \lambda_p, R_e)$.

549 Let $|h_{x,y^*}|^2$ be the channel power gain between x and y^* .
550 Therefore, when a PU-receiver exists, the received beacon
551 power (P_R) at x from y^* is given by $P_R = P_b |h_{x,y^*}|^2 g(r_1^*)$,
552 where $g(r_1^*)$ is the path-loss factor between x and y^* .

553 $P_{md}(x)$ can thus be written as

$$554 \quad P_{md}(x) = e^{-\pi \lambda_p R_e^2} + (1 - e^{-\pi \lambda_p R_e^2}) \times \Pr[R_b < P_{th}] \\ 555 \quad = e^{-\pi \lambda_p R_e^2} + (1 - e^{-\pi \lambda_p R_e^2}) \times \Pr \left[|h_{x,y^*}|^2 < \frac{P_{th}}{P_b g(r_1^*)} \right] \\ 556 \quad = e^{-\pi \lambda_p R_e^2} + (1 - e^{-\pi \lambda_p R_e^2}) \left(1 - e^{-\frac{P_{th}}{P_b}} \left(\frac{1 - e^{-\pi \lambda_p}}{1 - e^{-\pi \lambda_p R_e^2}} \right) \right. \\ 557 \quad \left. - \int_1^{R_e} \frac{2\pi \lambda_p t}{1 - e^{-\pi \lambda_p R_e^2}} e^{-\frac{P_{th}}{P_b t^{-\alpha}}} e^{-\pi \lambda_p t^2} dt \right), \quad (9)$$

558 and the integration in (9) can be performed numerically.

559 B. Co-Operative Spectrum Sensing

560 In this section, we analyze P_{md} when each CU employs
561 the CU selection schemes proposed in Section II-D. The total
562 P_{md} depends on both: 1) beacon misdetection, and 2) control
563 channel misdetection.

564 1) *Nearest Scheme*: Let the closest neighbour from CU $x \in$
565 Φ_s be denoted as x^* ($x^* \in \Phi_s$) with $x^* = \arg \min_{z \in \Phi_s} \|z - x\|$,
566 located at a distance \tilde{r}^* from x . Because the signals from x^*
567 with $\tilde{r}^* > R_e$ are neglected due to path loss, there may be
568 an occasion where a node x^* does not exist for cooperation.
569 This probability ρ_0 is obtained as $\rho_0 = e^{-\pi \lambda_s R_e^2}$ using the
570 void probability of a PPP. Let \tilde{r}_1^* be the distance from x to x^*
571 whenever $\tilde{r}^* < R_e$. Thus \tilde{r}_1^* is distributed as $TRal(\pi \lambda_s, R_e)$.

572 Node x^* senses the presence of primary receiver beacons,
573 and passes that information in the form of binary informa-
574 tion in a narrow band channel using another control signal.
575 Let $P_{b,s}$ be the power of this control signal, and $|h_{x,x^*}|^2$ be
576 the channel power gain between x and x^* . Therefore, if the
577 received control signal power ($P_{R,s}$) at x from x^* is given by
578 $P_{R,s} = P_{b,s} |h_{x,x^*}|^2 g(\tilde{r}_1^*)$, where $g(\tilde{r}_1^*)$ is the path loss gain
579 between x and x^* .

580 The probability of misdetecting the control signal transmit-
581 ted by x^* , $q_{s,i}$, is obtained as

$$582 \quad q_{s,i} = \Pr[P_{R,s} < P_{th}] = E_{\tilde{r}_1^*} \left[1 - e^{-\frac{P_{th}}{P_{b,s} g(\tilde{r}_1^*)}} \right]. \quad (10)$$

583 After performing the averaging with respect to \tilde{r}_1^* , the simpli-
584 fied expression for $q_{s,i}$ is

$$585 \quad q_{s,i} = 1 - e^{-\frac{P_{th}}{P_{b,s}} \left(\frac{1 - e^{-\pi\lambda_s}}{1 - e^{-\pi\lambda_s R_e^2}} \right)} \\ 586 \quad - \int_1^{R_e} \frac{2\pi\lambda_s t}{1 - e^{-\pi\lambda_s R_e^2}} e^{-\frac{P_{th}}{P_b t^{-\alpha}}} e^{-\pi\lambda_s t^2} dt \quad (11)$$

587 Let P_{md}^1 be the final misdetection probability of x when
588 cooperating with its closest neighbor. We will assume that x
589 uses an OR rule [11] where P_{md}^1 becomes the product of the
590 separate primary beacon and secondary control signal misde-
591 tecting probabilities. However, the probability that there is no
592 CU within R_e must be considered. P_{md}^1 is composed of the fol-
593 lowing events: (1) x^* does not exist, and x misdetects, (2) x^*
594 does exist, but both x^* and x misdetect the primary beacons,
595 and (3) x^* does exist, and detects the primary beacon cor-
596 rectly, but x misdetects both the primary system beacons and
597 the control signal from x^* . After combining these three events,
598 we can write P_{md}^1 as

$$599 \quad P_{md}^1 = P_{md}(x) \left(e^{-\pi\lambda_s R_e^2} + \left(1 - e^{-\pi\lambda_s R_e^2} \right) \right. \\ 600 \quad \left. \times \left(P_{md}(x) + (1 - P_{md}(x))q_{s,i} \right) \right). \quad (12)$$

601 We have used the fact that correct secondary control signal
602 reception due to double errors (x^* misdetects the primary bea-
603 cons but x detects a secondary control signal when it's not
604 present) are negligible. Moreover, spatial correlations have not
605 been taken into account in the derivation of (12).

606 2) *Multiple Random Scheme*: Let x_r ($x_r \in \Phi_s$) be any CU
607 within a cooperating distance of R_c from x , and r_r be the
608 distance from x to x_r . Using similar arguments as the derivation
609 of $\|x - y\|$, the distribution of r_r is shown to be distributed
610 according to $Lin(R_c)$.

611 Similar to the nearest scheme, whenever an x_r detects the
612 primary beacons, this information is sent via a control signal to
613 x . We assume that x can differentiate the control signals com-
614 ing from the M associated CUs, which can be easily achieved
615 via orthogonal codes serving as an identifier of each CU within
616 Φ_s . If $|h_{x,x_r}|^2$ and $g(r_r)$ are the small scale channel gain and
617 path loss gain between x_r and x , the received signal power
618 $P_{R,s}$ from x_r is given by $P_{R,s} = P_{b,s} |h_{x,x_r}|^2 g(r_r)$.

619 If $q_{s,i}$ is the probability of x misdetecting the control signal
620 from x_r , it is obtained as

$$621 \quad q_{s,i} = 1 - \frac{e^{-\frac{P_{th}}{P_{b,s}}}}{R_c^2} - \frac{2}{R_c^2} \int_1^{R_c} e^{-\frac{P_{th}}{P_{b,s} t^{-\alpha}}} t dt \\ 622 \quad = 1 - \frac{e^{-\frac{P_{th}}{P_{b,s}}}}{R_c^2} + \frac{2}{\alpha} E_{1-\frac{2}{\alpha}} \left(\frac{P_{th} R_c^\alpha}{P_{b,s}} \right). \quad (13)$$

623 Let P_{md}^2 be the final misdetection probability of $x \in \Phi_s$.
624 Although M is fixed beforehand, due to spatial randomness,
625 the available number of CUs may be less than M . Thus, P_{md}^2
626 is the sum of several probability components corresponding
627 to the number of cooperating nodes. Let q be the probabil-
628 ity of misdetection arising from a single cooperating node
629 (sum of the primary beacon misdetection probability by x_r

and the probability that the control signal of x_r is misdetected
by x when x_r correctly detects the primary beacons). It can
be written as $q = (P_{md}(x) + (1 - P_{md}(x))q_{s,i})$. Whenever a
given k ($k \leq M$) cooperating nodes are present, the final mis-
detecting probability of $x \in \Phi_s$ becomes $P_{md}(x)q^k$. As such,
 $P_{md}^2 = E_k[P_{md}(x)q^k]$, where $0 \leq k \leq M$. After averaging with
respect to k using (1), P_{md}^2 becomes

$$630 \quad P_{md}^2 = P_{md}(x) \left(e^{-\pi\lambda_s R_c^2 (1-q)} \frac{\Gamma(M, \pi\lambda_s R_c^2 q)}{\Gamma(M)} \right. \\ 631 \quad \left. + \left(1 - \frac{\Gamma(M, \pi\lambda_s R_c^2)}{\Gamma(M)} \right) q^M \right). \quad (14) \quad 632$$

633 3) *Best Received Power Scheme*: Let the neighbouring CU
634 of $x \in \Phi_s$ having the best instantaneous received signal
635 power be denoted as x_h . In order to evaluate the secondary
636 control signal misdetection probability ($q_{s,i}$), the Mapping
637 theorem [24] is used on the PPP Φ_s . Furthermore, for conven-
638 ience, we will use the path loss function $g(r_h) = r_h^{-\alpha}$ where
639 $r_h = \|x - x_h\|$ is the distance between x and x_h . Moreover,
640 we denote the channel gain between x and x_h as $|h_{x,x_h}|^2$. The
641 mapping procedure is as follows. With respect to $x \in \Phi_s$, the
642 process of CUs is homogeneous in \mathbb{R}^2 with it at the center.
643 It is shown that an inhomogeneous PPP $\Phi_{s,h}$ with intensity
644 $\lambda_{s,h}$, an exponential path loss with a path loss exponent of 1
645 and no fading generates the equivalent received power to that
646 from a homogeneous PPP, and exponential path loss with an
647 exponent α and Rayleigh fading [53], where $\lambda_{s,h}$ is written as
648 (see the Appendix)

$$649 \quad \lambda_{s,h} = \frac{2\pi}{\alpha} \lambda_s r_{s,h}^{\frac{2}{\alpha}-1} \Gamma\left(\frac{2}{\alpha} + 1\right), \quad 0 < r_{s,h} < \infty. \quad (15) \quad 650$$

651 Note that $r_{s,h}$ is a distance based metric of the PPP and not
652 any physical distance. In $\Phi_{s,h}$, the node having the smallest
653 distance metric from x is x_h . Thus, using (15), the PDF of the
654 distance metric to x_h (denoted by r_h^*) can be obtained as

$$655 \quad f_{r_h^*}(t) = \frac{2\pi}{\alpha} \lambda_s \Gamma\left(\frac{2}{\alpha} + 1\right) t^{\frac{2}{\alpha}-1} e^{-\pi\lambda_s \Gamma\left(\frac{2}{\alpha} + 1\right) t^{\frac{2}{\alpha}}}, \quad 0 < t < \infty. \quad (16) \quad 656$$

657 With these results, the received secondary control signal
658 power at x is written as $P_{R,s} = P_{b,s}(r_h^*)^{-1}$. Thus, $q_{s,i}$ is
659 obtained as

$$660 \quad q_{s,i} = \Pr\left[P_{b,s}(r_h^*)^{-1} < P_{th} \right] \\ 661 \quad = e^{-\pi\lambda_s \Gamma\left(\frac{2}{\alpha} + 1\right) \left(\frac{P_{b,s}}{P_{th}}\right)^{\frac{2}{\alpha}}}. \quad (17) \quad 662$$

663 The final misdetection probability of $\phi_{s,i}$ (P_{md}^3) is composed
664 of two components. First x and x_h may both misdetect the
665 primary beacon. Second, while x_h detects the primary beacon,
666 x may misdetect the control channel between x and x_h . Thus,
667 P_{md}^3 is obtained as

$$668 \quad P_{md}^3 = P_{md}(x) \left(P_{md}(x) + (1 - P_{md}(x))q_{s,i} \right). \quad (18) \quad 669$$

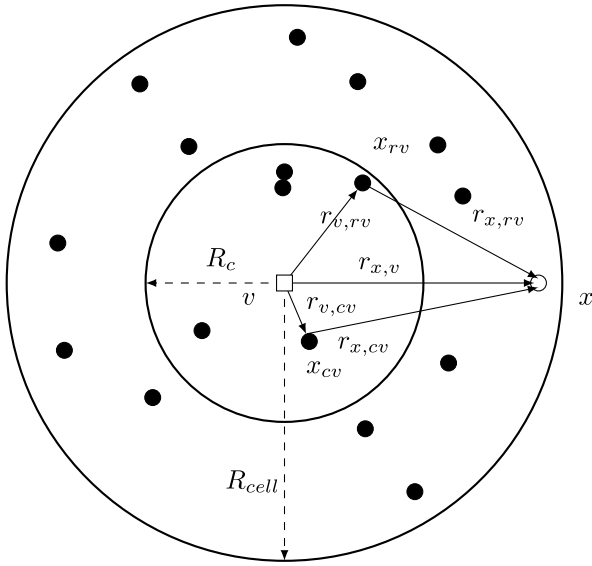


Fig. 2. The PU-transmitter v located at $(0,0)$ sends the beacon. The cell radius is denoted by R_{cell} , the cooperating radius is denoted by R_c , while the black dots denote the CUs. The CU x located at a distance $r_{x,v}$ from v can cooperate with either the closest CU to v (x_{cv}), or cooperate with a random CU within a distance of R_c from v (x_{rv}).

673 IV. P_{md} ANALYSIS FOR PU-TRANSMITTER BEACONS

674 This case is depicted in Fig. 2. In primary cellular networks
675 where the transmitter is a base station, and receivers are user
676 equipment, this approach provides wide benefits as base stations
677 are not power limited and avoids PU-receiver power
678 drain.

679 A. Local Primary Beacon Detection

680 Each CU ($x \in \Phi_s$) listens for the beacon of the PU-
681 transmitter (v) of its cell. Let R_{cell} be the cell radius, and
682 $P_{b,p}$ be power level of the beacon. Let $r_{x,v} = \|x - v\|$. This is
683 the distance between a fixed point and a random point from
684 Φ_s . The $r_{x,v}$ will be distributed as $Lin(R_{cell})$ (we assume that
685 $R_{cell} \ll R_e$). If $|h_{x,v}|^2$ and $g(r_{x,v})$ are the small scale chan-
686 nel gain and path loss gain between x and v , the received
687 beacon power at x (P_R) is given by $P_R = P_{b,p}|h_{x,v}|^2g(r_{x,v})$.
688 Whenever it falls below the threshold, the detection fails. Thus,
689 the probability of misdetection is given by

$$\begin{aligned}
 690 \quad P_{md}(x) &= \Pr\left[P_{b,p}|h_{x,v}|^2g(r_{x,v}) < P_{th}\right] \\
 691 \quad &= 1 - \frac{e^{-\frac{P_{th}}{P_{b,p}}}}{R_{cell}^2} - \frac{2}{R_{cell}^2} \int_1^{R_{cell}} e^{-\frac{P_{th}}{P_{b,p}t^{-\alpha}}} t dt \\
 692 \quad &= 1 - \frac{e^{-\frac{P_{th}}{P_{b,p}}}}{R_{cell}^2} + \frac{2}{\alpha} E_{1-\frac{2}{\alpha}}\left(\frac{P_{th}R_{cell}^\alpha}{P_{b,p}}\right). \quad (19)
 \end{aligned}$$

693 B. Co-Operative Sensing

694 For PU-transmitter emitted beacons, we will now analyze
695 the two additional schemes proposed.

696 1) *Nearest CU to PU-Transmitter Scheme*: Let x_{cv} be the
697 closest CU ($\in \Phi_s$) to v ($x_{cv} = \arg \min_{z \in \Phi_s} \|z - v\|$), with
698 $r_{v,cv} = \|v - x_{cv}\|$ and $r_{x,cv} = \|x - x_{cv}\|$. If $r_{v,cv} > R_{cell}$, a

cooperating node does not exist. The probability of this sce- 699
nario occurring (ρ_1) is given by $\rho_1 = e^{-\pi\lambda_s R_{cell}^2}$. Thus, the 700
variable $r_{v,cv}$ is distributed according to $TRal(\pi\lambda_s, R_{cell})$. This 701
distribution is obtained by removing x from Φ_s . This removal 702
does not significantly affect the statistics of Φ_s . 703

We now need to find the probability that x_{cv} misdetects 704
the PU-transmitter's beacon ($P_{md}(x_{cv})$) for this scenario. Let 705
 $|h_{v,cv}|^2$ and $g(r_{v,cv})$ be the small scale channel gain and path 706
loss gain between v and x_{cv} . The received beacon power at x_{cv} 707
($P_{R,s}$) is given by $P_{R,s} = P_{b,p}|h_{v,cv}|^2g(r_{v,cv})$. $P_{md}(x_{cv})$ is thus 708
obtained as 709

$$\begin{aligned}
 P_{md}(x_{cv}) &= 1 - e^{-\frac{P_{th}}{P_{b,p}} \left(\frac{1 - e^{-\pi\lambda_s}}{1 - e^{-\pi\lambda_s R_{cell}^2}} \right)} \\
 &\quad - \int_1^{R_{cell}} \frac{2\pi\lambda_s t}{1 - e^{-\pi\lambda_s R_{cell}^2}} e^{-\frac{P_{th}}{P_{b,p}t^{-\alpha}}} e^{-\pi\lambda_s t^2} dt. \quad (20) \quad 711
 \end{aligned}$$

We now derive the probability that $x \in \Phi_s$ misdetects sec- 712
ondary control signal from x_{cv} whenever it (x_{cv}) detects the 713
PU transmitter's beacon. The small scale channel gain and 714
path loss gain between x and x_{cv} are denoted by $|h_{x,cv}|^2$ and 715
 $g(r_{x,cv})$ respectively. The received power of the secondary con- 716
trol at $x \in \Phi_s$ is given by $P_{R,s} = P_{b,s}|h_{x,cv}|^2g(r_{x,cv})$, and the 717
probability of x misdetecting the control signal ($q_{s,i}$) is then 718
given by 719

$$\begin{aligned}
 q_{s,i} &= \Pr[P_{R,s} < P_{th}] \quad 720 \\
 &= E_{r_{x,cv}} \left[1 - e^{-\frac{P_{th}}{P_{b,s}g(r_{x,cv})}} \right]. \quad (21) \quad 721
 \end{aligned}$$

In order to evaluate this, the distribution of $r_{x,cv}$ is needed. 722
From the cosine rule, $r_{x,cv}$ can be written as $r_{x,cv} =$ 723
 $\sqrt{r_{v,cv}^2 + r_{x,v}^2 - 2r_{x,v}r_{v,cv} \cos \theta}$, where θ is a uniform between 724
 0 and 2π , with $f_\theta(x) = \frac{1}{2\pi}$, $0 \leq x < 2\pi$. Furthermore, for 725
mathematical convenience, we will take $g(r_{x,cv}) = r_{x,cv}^{-\alpha}$. Thus, 726
 $q_{s,i}$ becomes 727

$$\begin{aligned}
 q_{s,i} &= 1 - \int_0^{R_{cell}} \int_0^{R_{cell}} \int_0^{2\pi} e^{-\frac{P_{th}}{P_{b,s}(r_{v,cv}^2 + r_{x,v}^2 - 2r_{x,v}r_{v,cv} \cos \theta)^{-\frac{\alpha}{2}}}} \\
 &\quad \times \frac{2\lambda_s r_{x,v} r_{v,cv}}{R_{cell}^2 (1 - e^{-\pi\lambda_s R_{cell}^2})} e^{-\pi\lambda_s r_{v,cv}^2} d\theta dr_{v,cv} dr_{x,v}. \quad (22) \quad 729
 \end{aligned}$$

Let P_{md}^4 be the overall misdetection probability of $x \in \Phi_s$. 730
Similar to the previous analysis, it is necessary to con- 731
sider probability of no cooperating node (ρ_1). Thus, P_{md}^4 is 732
composed of three events: (1) x misdetects beacon and no 733
cooperating node exists, (2) x and x_{cv} both misdetect beacon, 734
and (3) x misdetects the beacon and x_{cv} detects it but x mis- 735
detects the control signal from x_{cv} . Considering these three 736
events, we can write 737

$$\begin{aligned}
 P_{md}^4 &= P_{md}(x) \left(e^{-\pi\lambda_s R_{cell}^2} + \left(1 - e^{-\pi\lambda_s R_{cell}^2} \right) \right. \\
 &\quad \left. \times \left(P_{md}(x_{cv}) + (1 - P_{md}(x_{cv}))q_{s,i} \right) \right). \quad (23) \quad 739
 \end{aligned}$$

2) *Random CU to PU-Transmitter Scheme*: Let the ran- 740
domly selected CU be x_{rv} , its distance from v be $r_{v,rv}$, and its 741
distance from x be $r_{x,rv}$. We assume that $R_c < \min(R_{cell}, R_e)$. 742

If no such CU exists within a distance of R_c of v , no cooperation occurs. The probability of it is $\rho_2 = e^{-\pi\lambda_s R_c^2}$.

The probability that x_{rv} misdetects the beacon from v is obtained next. We denote this probability as $P_{md}(x_{rv})$, and the small scale channel gain and path loss gain between v and x_{rv} respectively as $|h_{v,rv}|^2$ and $g(r_{v,rv})$. The received beacon power at x_{rv} ($P_{R,s}$) is given by $P_{R,s} = P_{b,p}|h_{v,rv}|^2 g(r_{v,rv})$. We can now write $P_{md}(x_{rv})$ as

$$\begin{aligned} P_{md}(x_{rv}) &= \Pr\left[P_{b,p}|h_{v,rv}|^2 g(r_{v,rv}) < P_{th}\right] \\ &= 1 - \frac{e^{-\frac{P_{th}}{P_{b,p}}}}{R_c^2} - \frac{2}{R_c^2} \int_1^{R_c} e^{-\frac{P_{th}}{P_{b,p}t^{-\alpha}}} t dt \\ &= 1 - \frac{e^{-\frac{P_{th}}{P_{b,p}}}}{R_c^2} + \frac{2}{\alpha} E_{1-\frac{2}{\alpha}}\left(\frac{P_{th}R_c^\alpha}{P_{b,p}}\right). \end{aligned} \quad (24)$$

We will now derive the probability that x misdetects the secondary control signal from x_{rv} (denoted by $q_{s,i}$), whenever a secondary control signal is transmitted. The small scale channel gain and path loss gain between x and x_{rv} are respectively denoted as $|h_{x,rv}|^2$ and $g(r_{x,rv})$. Similar to the previous scheme, we will use $g(r_{x,rv}) = r_{x,rv}^{-\alpha}$ for mathematical convenience. Using the cosine rule, $r_{x,rv}$ is written as $r_{x,rv} = \sqrt{r_{v,rv}^2 + r_{x,v}^2 - 2r_{x,v}r_{v,rv} \cos \theta}$. Thus, $q_{s,i}$ is written as

$$\begin{aligned} q_{s,i} &= \Pr\left[P_{b,s}|h_{x,rv}|^2 r_{x,rv}^{-\alpha} < P_{th}\right] \\ &= 1 - \int_0^{R_{cell}} \int_0^{R_c} \int_0^{2\pi} e^{-\frac{P_{th}}{P_{b,s}(r_{v,rv}^2 + r_{x,v}^2 - 2r_{x,v}r_{v,rv} \cos \theta)^{-\frac{\alpha}{2}}}} \\ &\quad \times \frac{2}{\pi R_c^2 R_{cell}^2} r_{x,v} r_{v,rv} d\theta dr_{v,rv} dr_{x,v}. \end{aligned} \quad (25)$$

The final misdetection probability of x (denoted as P_{md}^5) is comprised of 3 terms as the previous scheme (Nearest CU to PU-transmitter scheme). Thus, P_{md}^5 is obtained as

$$\begin{aligned} P_{md}^5 &= P_{md}(x) \left(e^{-\pi\lambda_s R_c^2} + \left(1 - e^{-\pi\lambda_s R_c^2}\right) \right. \\ &\quad \left. \times \left(P_{md}(x_{rv}) + (1 - P_{md}(x_{rv}))q_{s,i} \right) \right). \end{aligned} \quad (26)$$

This scheme can be generalized where a cooperates with up to M CUs within a distance of R_c from the PU-transmitter.

V. P_f ANALYSIS

For completeness, we will conduct an analysis of the false alarm probability P_f . First, we will analyze P_f for the different local detection schemes for PU-receiver and PU-transmitter beacons.

A. P_f for Local Detection Schemes

1) *Aggregating Beacon Power*: A false alarm occurs when the CU detects the presence of a beacon when none are present. In this scenario, the received power is purely composed of noise. Thus

$$P_R = w, \quad (27)$$

where $w = \mathcal{N}(0, \sigma^2)$, and σ^2 is the noise variance (it should be noted that because a narrowband channel is used for beacons and control signals, σ^2 is very small). Let $P_f(x)$ be the probability of falsely detecting PU beacons by the CU $x \in \Phi_s$ in its local detection phase. P_f can be written as $P_f = \Pr[P_R > P_{th}]$. As such

$$P_f(x) = Q\left(\frac{P_{th}}{\sigma}\right), \quad (28)$$

and $Q(\cdot)$ is the Q function.

2) *Separately Sensing Primary Beacons*: When separately detecting primary beacons, a false alarm can occur even if a single stream from a PU is detected in error. Thus, we have

$$P_f(x) = \mathbb{E}\left[1 - (\Pr[P_R < P_{th}])^N\right] = \mathbb{E}\left[1 - \left(1 - Q\left(\frac{P_{th}}{\sigma}\right)\right)^N\right]. \quad (29)$$

After averaging with respect to N ,

$$P_f(x) = 1 - e^{-\pi R_c^2 \lambda_p Q\left(\frac{P_{th}}{\sigma}\right)}. \quad (30)$$

3) *Closest PU-Receiver Selection*: $P_f(x)$ for this scenario is identical to (28), and $P_f(x) = Q\left(\frac{P_{th}}{\sigma}\right)$.

4) *PU-Transmitter Beacons*: As each CU ($x \in \Phi_s$) listens to the beacon of the primary transmitter of its own cell, $P_f(x)$ is simply written similar to (28) as $P_f(x) = Q\left(\frac{P_{th}}{\sigma}\right)$.

B. P_f After Co-Operation

Using the local false alarm probabilities derived above, we now derive the final false alarm probability after co-operation for the different schemes.

1) *PU-Receiver Beacons (Nearest Scheme)*: For this scheme, false alarm occurs even if one of the following cases occur: 1) x falsely detect beacons, 2) x properly detects beacons, the nearest neighbour x^* properly detects, but x improperly detects the control channel, and 3) x properly detects beacons, the nearest neighbour x^* falsely detects, and x detects the control channel. After combining these events, we can write P_f^1 as

$$P_f^1 = P_f(x) + (1 - P_f(x))((1 - P_f(x))P_f(x) + P_f(x)(1 - q_{s,i})). \quad (31)$$

It should be noted that the probability of falsely detecting the control channel also follows (28), and that $q_{s,i}$ follows (11).

2) *PU-Receiver Beacons (Multiple Random Scheme)*: In this scheme, even a single false alarm from one of the co-operating nodes triggers a false alarm after combination. Let p be the probability that there is **no** false alarm from a co-operating node x_r . p can be written as $p = (1 - P_f(x))(1 - P_f(x)) + P_f(x)q_{s,i}$. The final false alarm probability can thus be written as $P_f^2 = 1 - (1 - P_f(x))p^k$, for a given $k(\leq M)$. Averaging with respect to k (1) results in

$$\begin{aligned} P_f^2 &= 1 - (1 - P_f(x)) \left(e^{-\pi\lambda_s R_c^2 (1-p)} \frac{\Gamma(M, \pi\lambda_s R_c^2 p)}{\Gamma(M)} \right. \\ &\quad \left. + \left(1 - \frac{\Gamma(M, \pi\lambda_s R_c^2)}{\Gamma(M)}\right) p^M \right). \end{aligned} \quad (32)$$

3) *PU-Receiver Beacons (Best Received Power Scheme)*:
The final false alarm probability P_f for this scheme follows (31) with $q_{s,i}$ following (17).

4) *PU-Transmitter Beacons (Nearest CU to PU-Transmitter Scheme)*: The final P_f for this scheme follows (31) with $q_{s,i}$ following (22).

5) *PU-Transmitter Beacons (Random CU to PU-Transmitter Scheme)*: The final P_f for this scheme also follows (31) with $q_{s,i}$ following (25).

VI. PRIMARY SYSTEM PERFORMANCE

The misdetection of beacons by a set of CUs will cause interference, which will degrade the received SINR, $\gamma_{p,y}$, at PU-receiver $y \in \Phi_p$. Thus, let I be the aggregate interference from the CUs, $P_{R,p}$ be the received primary signal power at $y \in \Phi_p$, and σ_n^2 be the noise power spectral density at the PU-receiver. We assume that different PU-transmitters use orthogonal codes, and do not pose significant interference to PU-receivers within other cells. $P_{R,p}$ is written as $P_{R,p} = P_p |h_{v,y}|^2 g(r_{v,y})$, where P_p , $|h_{v,y}|$ and $r_{v,y}$ are respectively the PU transmit power, channel power gain and distance between the PU-transmitter v and y . We can thus write the SINR as $\gamma_y = \frac{P_{R,p}}{I + \sigma_n^2}$. An outage occurs whenever $\gamma_y < \gamma_{th}$ where γ_{th} is a threshold. Note that we are more interested in the SINR falling below a threshold for the primary signals as opposed to the received signal falling below a threshold used for beacon detection. The primary signals would be transmitting data whereas the beacon signals only indicate the channel occupation for which the received signal level was sufficient. Thus, the outage probability of γ_y may be written as

$$P_{Out,y} = \Pr[\gamma_y < \gamma_{th}].$$

We can write

$$\begin{aligned} P_{Out,y/I,r_{v,y}}(x) &= \Pr\left[\frac{P_p |h_{v,y}|^2 g(r_{v,y})}{I + \sigma_n^2} \leq \gamma_{th}\right] \\ &= \Pr\left[|h_{v,y}|^2 \leq \frac{\gamma_{th}(I + \sigma_n^2)}{P_p g(r_{v,y})}\right] \\ &= 1 - e^{-\left(\frac{\gamma_{th}(I + \sigma_n^2)}{P_p g(r_{v,y})}\right)}. \end{aligned}$$

$P_{Out,y/I,r_{v,y}}(x)$ can be further averaged with respect to I as

$$\begin{aligned} P_{Out,j/r_{v,y}}(x) &= 1 - e^{-\left(\frac{\gamma_{th}\sigma_n^2}{P_p g(r_{v,y})}\right)} E_I\left[e^{-I\left(\frac{\gamma_{th}}{P_p g(r_{v,y})}\right)}\right] \\ &= 1 - e^{-\left(\frac{\gamma_{th}\sigma_n^2}{P_p g(r_{v,y})}\right)} M_I\left(\frac{\gamma_{th}}{P_p g(r_{v,y})}\right). \end{aligned} \quad (33)$$

Equation (33) provides the outage probability of node y given $r_{v,y}$ is known. However, if averaging over all PU-receivers is needed, we need the PDF of $r_{v,y}$, which is the distance from a fixed point to a random point from Φ_p , which can be shown

to be $Lin(R_{cell})$. Thus, the average outage can be expressed as

$$\begin{aligned} P_{Out,y}(x) &= 1 - \frac{e^{-\left(\frac{\gamma_{th}\sigma_n^2}{P_p}\right)}}{R_{cell}^2} M_I\left(\frac{\gamma_{th}}{P_p}\right) \\ &\quad - \int_1^{R_{cell}} 2 \frac{t}{R_{cell}^2} e^{-\left(\frac{\gamma_{th}\sigma_n^2}{P_p t^{-\alpha}}\right)} M_I\left(\frac{\gamma_{th}}{P_p t^{-\alpha}}\right) dt. \end{aligned} \quad (34)$$

To evaluate this, the MGF of the aggregate interference at y ($M_I(s)$) needs to be obtained. However, the exact expressions for interference is a function of each individual cooperating scheme, and thus complex. But, for completion, we suggest the following approximate approach. $M_I(s)$ is written as $M_I(s) = E[e^{-sI}]$. Let $r_{x,y} = \|x - y\|$ for any interfering CU $x \in \Phi_s$, which is distributed as $Lin(R_e)$. Note that similar to Section III, we do not consider the interference from x whenever $r_{x,y} > R_e$. When $P_{md}(x)$ is the final misdetection probability of $x \in \Phi_s$ with CBS, the Coloring theorem [24] suggests that the intensity of the interfering CUs is $P_{md}(x)\lambda_s$. $M_I(s)$ is thus obtained as

$$M_I(s) = e^{\pi R_e^2 P_{md}(x)\lambda_s (M_{I_x}(s) - 1)}, \quad (35)$$

where $M_{I_x}(s)$ is the MGF of the interference from x . It is given by $M_{I_x}(s) = E[e^{-sP_s |h_{x,y}|^2 g(r_{x,y})}]$, where P_s is the CU transmit power, and $|h_{x,y}|^2$ and $g(r_{x,y})$ are respectively the small scale channel gain and the path loss gain between y and x . $M_{I_x}(s)$ is derived as

$$M_{I_x}(s) = \frac{1}{R_e^2} \left(\sum_{k=0}^{\infty} (-sP_s)^k + \sum_{l=0}^{\infty} 2(-sP_s)^l \frac{R_e^{2-\alpha l} - 1}{2 - \alpha l} \right). \quad (36)$$

VII. NUMERICAL RESULTS

We will provide numerical results on the total misdetection and false alarm probabilities for the different cooperation and local primary beacon detection schemes. We used MATLAB for the simulation, with 10^4 topologies, and 10^4 transmissions for each topology; thus 10^8 simulations for each plot point. Note that because simulation results match with the theoretical results, we have not used separate marker styles.

A. Beacons Emitted by PU-Receiver Nodes

We will first investigate the case of PU-receiver beacons. The parameters are $R_e = 1500$, $R_c = 500$, $\alpha = 3$, $P_{b,s} = -40$ dBm, and $P_b = -50$ dBm. P_b has been set 10 dB lower than $P_{b,s}$ because it makes sense that the energy of a PU-receiver node should not be used excessively for beacon signaling. Moreover, P_{th} is chosen as -110 dBm, which is the minimum signal reception thresholds for several mobile standards [54].

Fig. 3 plots the total misdetection probability P_{md} (eqs. (14), (18), and (12)) and the false alarm probability P_{md} with respect to the CU detection threshold (P_{th}). While the performance improvement due to CBS is slight for higher P_{th} , it is significant when P_{th} is small. For example, when $P_{th} = -120$ dBm and using multiple random cooperation with 10 nodes, P_{md} decreases by a 10^4 fold. This decrease is even higher for best received power cooperation when separately

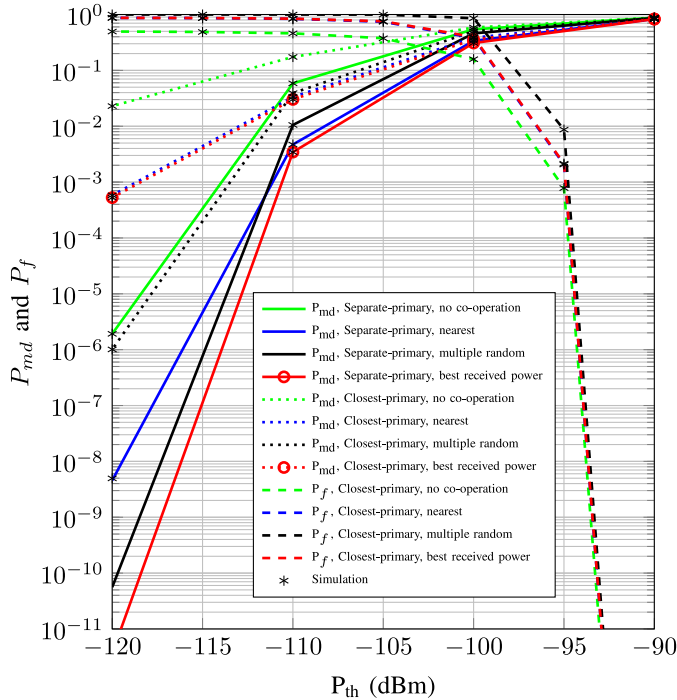


Fig. 3. P_{md} and P_f for PU-receiver beacons as a function of P_{th} for different cooperation schemes. $\lambda_p = 0.0001$, $\sigma^2 = 10^{-10}$, $\lambda_s = 0.0001$, and $M = 10$.

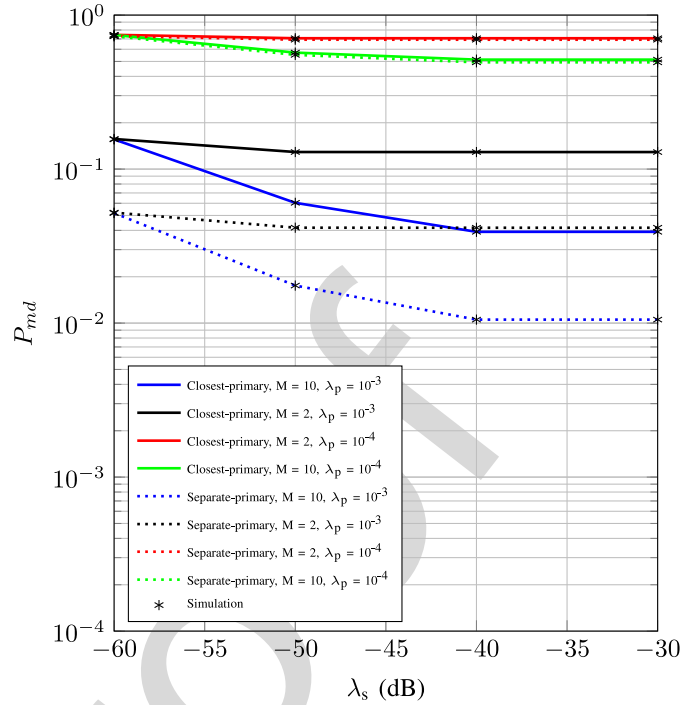


Fig. 4. P_{md} for PU-receiver beacons as a function of CU receiver density λ_s for multiple random cooperation. $P_{th} = -110$ dBm.

917 detecting all PU-receiver beacons. Furthermore, the latter
 918 performs better than sensing the beacon from the closest PU-
 919 receiver. However, as mentioned before, this comes at the cost
 920 of additional complexity and resources. It is also interesting
 921 to note that while the nearest scheme performs better than
 922 the multiple random scheme when P_{th} is higher, the con-
 923 verse is true for lower P_{th} . Moreover, while the best received
 924 power cooperation scheme always has better performance
 925 than the nearest scheme, the difference is only slight when
 926 detecting the closest PU-receiver's beacon. Contrary to the
 927 misdetection probability, the false alarm probability is very
 928 high for low P_{th} values and drops sharply as P_{th} increases.
 929 As expected, co-operation slightly increases the false alarm
 930 probability. The multiple random scheme with PU-receiver
 931 beacons has the worst performance because this scheme takes
 932 input from multiple CUs; even a single false alarm from
 933 one CU makes the final decision a false alarm. Moreover,
 934 the nearest and best received power co-operation schemes
 935 show almost identical performance with respect to the false
 936 alarm probability. It should also be noted that as unlicensed
 937 users, CUs should err in the side of false alarm rather than
 938 misdetection.

939 The behaviour of P_{md} for the multiple random scheme
 940 (eq. (14)) is investigated in Fig. 4 under different values of
 941 M and primary node density λ_p . For both separate and closest
 942 methods of primary beacon detection, the misdetection prob-
 943 ability approaches 1 when λ_p is low. Increasing the number
 944 of cooperating nodes M does not help significantly. However,
 945 when λ_p increases to 10^{-3} , increasing M has some effect.
 946 Furthermore, the performance gap between these two methods
 947 becomes apparent. Moreover, all curves flatten out indicating
 948 that the effect of λ_s becomes negligible beyond -40 dB.

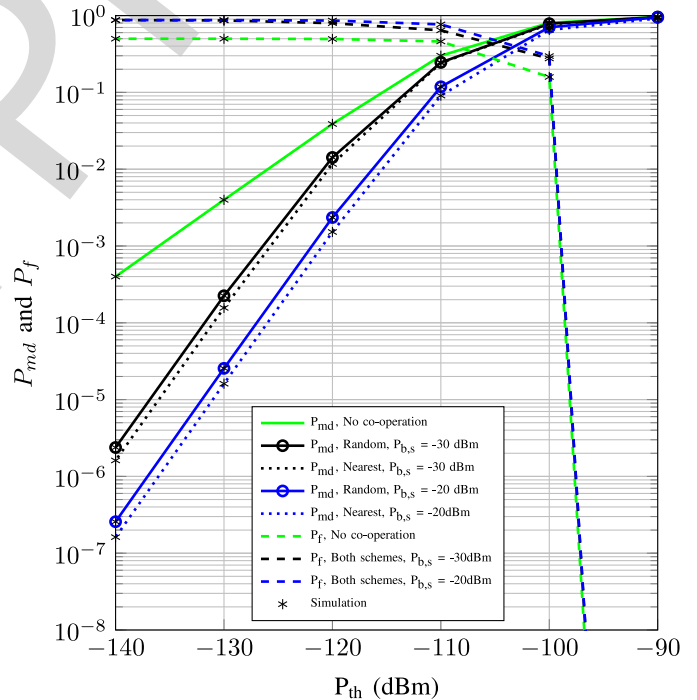


Fig. 5. P_{md} and P_f as a function of P_{th} for PU-transmitter beacons. $\lambda_s = 0.0001$, $R_{cell} = 1000$, $R_c = 500$, $\sigma^2 = 10^{-10}$ and $P_{b,p} = -20$ dBm.

B. Beacons Emitted by PU-Transmitter Nodes

949 We now focus on nearest CU-to-PU and random CU-to-
 950 PU schemes (eqs. (23) and (26)). Parameter values of $\alpha = 3$
 951 and $P_{b,p} = -20$ dBm are used. The latter reflects the fact
 952 that the PU-transmitters can manage high power levels. Fig. 5
 953 shows how P_{md} and P_f of the two CBS schemes varies with
 954

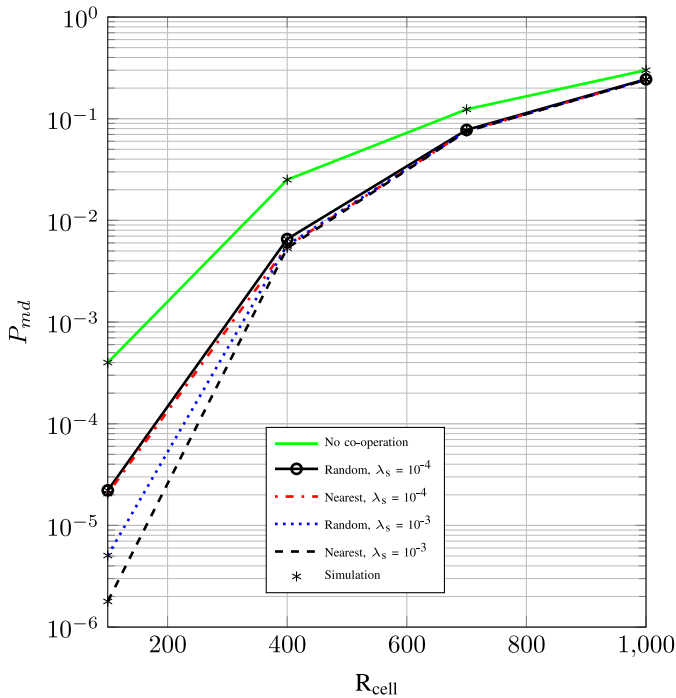


Fig. 6. P_{md} for PU-transmitter beacons as a function of R_{cell} . $R_c = 100$, $P_{b,p} = -20$ dBm, $P_{b,s} = -30$ dBm, and $P_{th} = -110$ dBm.

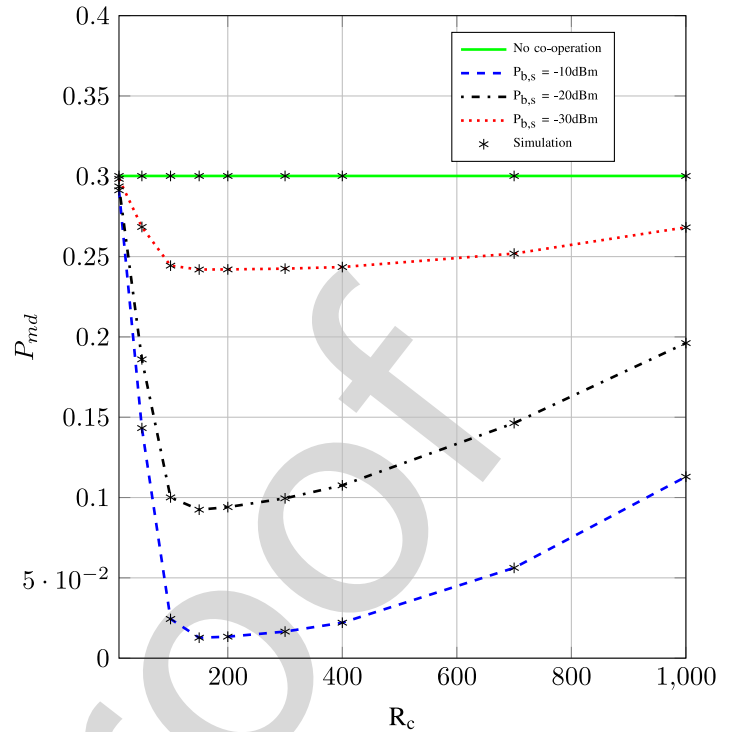


Fig. 7. P_{md} for PU-transmitter beacons as a function of R_c for random cooperation. $R_{cell} = 1000$, $P_{b,p} = -20$ dBm, and $P_{th} = -110$ dBm.

the detection threshold P_{th} . The effect of cooperation is more pronounced for low P_{th} values in terms of misdetection. The impact of the control channel is also seen. For example, a 10 dB increase in the control power $P_{b,s}$ results in order of magnitude reduction of misdetection. For both $P_{b,s}$ values, the nearest CU to PU-transmitter scheme has a slightly lower misdetection probability compared with the random CU to PU-transmitter scheme. In terms of false alarm, cooperation slightly increases P_f , and both co-operation schemes show very similar performance. When P_{th} increases beyond -100 dBm, there is a sudden drop in the false alarm probability. Furthermore, as expected, when the control channel power increases, the false alarm probability is slightly higher as erroneous information is more readily received from co-operating devices.

In Fig. 6, we study the impact of the cell size; P_{md} (eqs. (23) and (26)) versus cell radius R_{cell} is plotted. The most important insight from this graph is that the effect of cooperation decreases as cell radius increases when other parameters are kept constant, and that both CBS schemes converge in performance. This is due to a high R_{cell} outweighing the effect from other parameters, and the overall performance gain diminishing. With a high R_{cell} and a low cooperation radius R_c , the distance from the given CU to its cooperating node is high irrespective of the cooperation scheme causing similar secondary control channel misdetection probabilities. As expected, increasing the CU node spatial density λ_s decreases P_{md} . This is especially important for the nearest CU to PU-transmitter scheme. For the random CU to PU-transmitter scheme, increasing λ_s ensures that there is a CU available for cooperation within R_c . The effect of increasing λ_s are mainly seen for lower cell radius values. Furthermore, the nearest CU

to PU-transmitter scheme shows a slightly better performance than the random CU to PU-transmitter scheme for both λ_s values. However, the performance increase is higher when $\lambda_s = 10^{-3}$.

The effect of the cooperation radius R_c on the misdetection for the random CU to PU-transmitter scheme (26) is investigated in Fig. 7 for various levels of control signal power, $P_{b,s}$. A best-case cooperation radius can be observed, which ensures the lowest misdetection probability. When the cooperation radius approaches 0, random cooperation approaches convergence no cooperation as expected. However, as R_c increases, the misdetection probability drops steeply to the best-case value. Furthermore, it is observed that the steepness of this reduction increases with the control signal power. Subsequent increases in R_c up to R_{cell} only result in a gradual increase in misdetection.

VIII. CONCLUSION

This paper investigated the overall misdetection and false alarm probabilities of an interweave CU using several cooperative beacon sensing strategies. We captured the spatial randomness of PU and CU nodes via independent PPPs. The propagation effects included path loss and Rayleigh fading. Moreover, beacons emitted by both PU-receivers and PU-transmitters were considered. For the former, when sensing beacons emitted by the closest PU-receiver, multiple random CBS performs better when the reception threshold P_{th} is lower; e.g., misdetection decreases by 10^4 fold for thresholds as low as -120 dBm. However, the best received power scheme works slightly better for higher P_{th} . Moreover, the

1016 spatial density of PU-receiver nodes varies inversely with
 1017 detection performance. Furthermore, the best received power
 1018 scheme outperforms the nearest and multiple random cooper-
 1019 ation schemes significantly when CUs sense primary beacons
 1020 separately. When PU-transmitters send the beacons, a 10 dB
 1021 increase in $P_{b,s}$ decreases the misdetection probability by 10
 1022 fold for both cooperation schemes. Furthermore, the effect
 1023 of cooperation decreases for higher cell radii, and there
 1024 exists a best case cooperation distance R_c which provides the
 1025 lowest misdetection probability for random cooperation. For
 1026 PU-transmitter beacons, nearest cooperation provides slightly
 1027 better results than random cooperation. In addition, it was seen
 1028 that co-operation slightly increases the false alarm probability.
 1029 Future research ideas extending this work include considering
 1030 spatial and temporal correlation, considering other detection
 1031 rules at the CU, and investigating the energy efficiency of
 1032 cooperation schemes for CR networks.

APPENDIX

PROOF OF EQUATION (15)

1033 The density function associated with a PPP in \mathbb{R}^2 can be
 1034 transformed to polar coordinates using the Mapping theo-
 1035 rem [24] (This is used to convert the 2-D PPP to a 1-D PPP).
 1036 Thus, the density of the 1-D PPP of CUs with respect to x
 1037 ($\lambda_{r,1}(r)$) can be obtained as

$$1040 \lambda_{r,1}(r) = \int_0^{2\pi} \lambda_s r_{r,1} d\theta = 2\pi \lambda_s r_{r,1}, \quad 0 < r_{r,1} < \infty. \quad (37)$$

1041 The received power at a CU x from a cooperating CU at dis-
 1042 tance r is given by $Pr^{-\alpha}|h|^2$, where P is the transmit power
 1043 and $|h|^2$ is the channel gain between a cooperating CU and x .
 1044 Our first objective is to use the Mapping theorem to obtain a
 1045 new equivalent PPP which generates a received power identi-
 1046 cal to what is generated by the above PPP with intensity $\lambda_{r,1}$,
 1047 but with a path loss exponent of 1. The intensity function of
 1048 the new PPP $\lambda_{r,2}(r)$ is derived as [53]

$$1049 \lambda_{r,2}(r) = \frac{2\pi \lambda_s r_{r,2}^{\frac{2}{\alpha}-1}}{\alpha}, \quad 0 < r_{r,2} < \infty. \quad (38)$$

1050 In the next step, we use the Marking theorem [24] and the
 1051 Mapping theorem to obtain a new PPP which generates the
 1052 identical received power, but with a path loss exponent of 1
 1053 and no fading. The intensity function of the new PPP $\lambda_{s,hp}(r)$
 1054 can be derived as [53]

$$1055 \lambda_{s,h}(r) = E_{|h|^2} \left[|h|^2 \lambda_{r,2} \left(r_{s,h} |h|^2 \right) \right], \quad 0 < r_{s,h} < \infty. \quad (39)$$

1056 When the fading is modelled as Rayleigh, (39) can be
 1057 simplified as

$$1058 \lambda_{s,h}(r) = \frac{2\pi}{\alpha} \lambda_s r_{s,h}^{\frac{2}{\alpha}-1} E_{|h|^2} \left[\left(|h|^2 \right)^{\frac{2}{\alpha}} \right], \quad 0 < r < \infty$$

$$1059 = \frac{2\pi}{\alpha} \lambda_s r_{s,h}^{\frac{2}{\alpha}-1} \Gamma \left(\frac{2}{\alpha} + 1 \right), \quad 0 < r_{s,h} < \infty, \quad (40)$$

1060 which is equation (15). It should be noted that the limits of r
 1061 do not change because the CUs are distributed in a 2-D field.

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